

**EE 109 Homework 3**  
**Redekopp**

Name:     Solutions    

Due: \_\_\_\_\_ Score: \_\_\_\_\_

**Show work to get full credit. Remember, use on only one side of the paper and staple them together. Only use a calculator to CHECK your work, not to DO your work.**

1.a  $F = (A' + B')(C)$

ABC	A'	B'	A' + B'	C	F
000	1	1	1	0	0
001	1	1	1	1	1
010	1	0	1	0	0
011	1	0	1	1	1
100	0	1	1	0	0
101	0	1	1	1	1
110	0	0	0	0	0
111	0	0	0	1	0

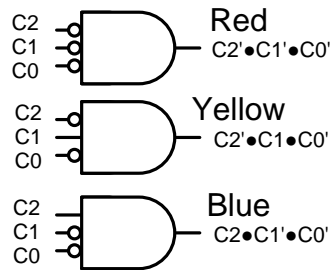
1.b  $G = WX' + XY' + W'$

WXY	X'	W • X'	Y'	X • Y'	W'	G
000	1	0	1	0	1	1
001	1	0	0	0	1	1
010	0	0	1	1	1	1
011	0	0	0	0	1	1
100	1	1	1	0	0	1
101	1	1	0	0	0	1
110	0	0	1	1	0	1
111	0	0	0	0	0	0

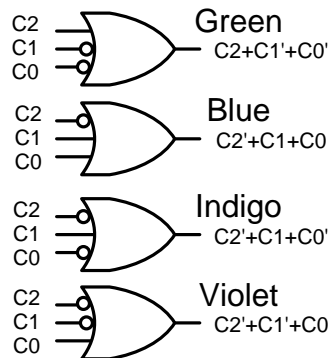
2. [BB] (14 pts.) Implement the following functions

Meaning	Color Bit Combinations		
	C2	C1	C0
Red	0	0	0
Orange	0	0	1
Yellow	0	1	0
Green	0	1	1
Blue	1	0	0
Indigo	1	0	1
Violet	1	1	0
Unused	1	1	1

- a. Build appropriate single gate (+ inverters) checker/decoders that outputs a logic '1' for each of the three primary colors: **Red, Blue, Yellow** (i.e. one decoder to check for each color)

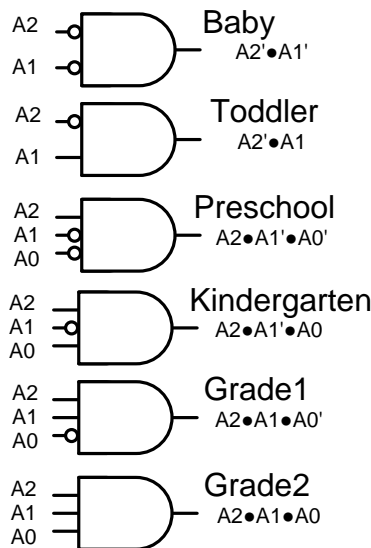


- b. Build an appropriate single gate (+inverters) checker/decoder that outputs a logic '0' for each color: **Green, Blue, Indigo, Violet**.



3. **[BB]** (12 pts.) Given a child's age [0-7 years old], design logic to produce the following outputs: {**Baby, Toddler, Preschool, Kinder, Grade1, Grade2**} which have the following age correspondence/mapping. Use only 1 AND gate (and any number of inverters) per output.

Meaning	Age Combinations		
	A2	A1	A0
Baby	0	0	0
	0	0	1
Toddler	0	1	0
	0	1	1
Preschool	1	0	0
Kindergarten	1	0	1
Grade1	1	1	0
Grade2	1	1	1



4.  $H = \Sigma_{A[3:0]}(0,3,6,9,12,15) =$   
 $A3'A2'A1'A0' + A3'A2'A1A0 + A3'A2A1A0' + A3A2'A1'A0 + A3A2A1'A0' +$   
 $A3A2A1A0$
5.  $G = \Sigma_{ABCD}(2,3,6,13,15)$   
 $= A'B'CD' + A'B'CD + A'BCD' + ABC'D + ABCD$
6.  $G = A'B'CD' + A'B'CD + A'BCD' + ABC'D + ABCD = m2 + m3 + m6 + m13 + m15$   
 Notice m2 and m3 differ by 1 variable, m2 and m6 differ by 1 variable, and m13 and m15 differ by 1 variable... We can replicate m2 (because T3:  $X+X = X$ ) to use m2 to help us with m3 and m6

$$\begin{aligned}
&= (m2+m3) + (m2+m6) + (m13+m15) \\
&= (A'B'CD' + A'B'CD) + (A'B'CD' + A'BCD') + (ABC'D + ABCD) \\
&= A'B'C(D'+D) + A'CD'(B'+B) + ABD(C'+C) \\
&= A'B'C + A'CD' + ABD
\end{aligned}$$

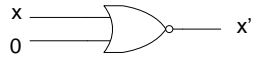
7. a.

$$\begin{aligned}
F &= X + [(W'Y'Z)(W + (X'(Y+Z)))] \\
&= X + [W'Y'ZW + W'Y'ZX'(Y+Z)] && \text{T8 Distributing} \\
&= X + [W'Y'ZW + W'Y'ZX'Y + W'Y'ZX'Z] && \text{T8 Distributing} \\
&= X + [0 \cdot Y'Z + W'ZX' \cdot 0 + W'Y'ZX'] && \text{T5' (W'W = 0 and Y'Y = 0) \& T3'} \\
&= X + [0 + 0 + W'X'Y'Z] && \text{T2'} \\
&= X + [W'X'Y'Z] && \text{T1} \\
&= (X+W')(X+X')(X+Y')(X+Z) && \text{T8' (Distributing + over \cdot)} \\
&= (X+W')(X+Y')(X+Z) && \text{T5, T1'}
\end{aligned}$$

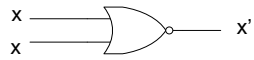
b

$$\begin{aligned}
G &= \overline{(w+x+y)} + \overline{wxy} + \overline{(wz+xy)} \cdot \overline{(x(\overline{w+z}))} \\
&= \overline{w} \cdot \overline{(x+y)} + \overline{wxy} + \overline{(wz+xy)} \cdot \overline{(x + (\overline{w+z}))} && \text{-- DeMorgan's theorem} \\
&= \overline{wxy} + \overline{wxy} + \overline{(wz+xy)} \cdot \overline{(x + \overline{wz})} \\
&= \overline{wxy} + \overline{wxy} + \overline{wxz} + \overline{xy} + \overline{wxz} + \overline{xy} \\
&= (w + \overline{w})xy + \overline{wxz} + 0 + \overline{wz} + \overline{wxyz} && \text{-- } x\overline{x} = 0 \\
&= \overline{xy} + \overline{wxz} + \overline{wz} + \overline{wxyz} \\
&= \overline{xy} + \overline{wz}(x+1) + \overline{wxyz} && \text{-- } 1 + \overline{x} = 1 \\
&= \overline{xy} + \overline{wz} + \overline{wxyz} \\
&= \overline{xy} + \overline{wz}(1 + \overline{xy}) && \text{-- } 1 + x\overline{y} = 1 \\
&= \overline{xy} + \overline{wz}
\end{aligned}$$

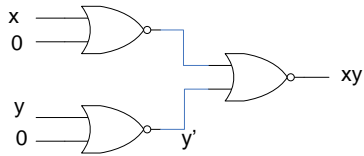
8.a  $x' = x \text{ NOR } 0$



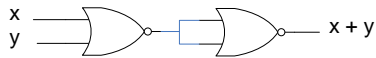
or inversion can be achieved as



8.b According to DeMorgan's theorem,  $x \cdot y = \overline{\overline{x + y}}$



8.c  $x + y = \overline{\overline{x + y}}$



8.d Yes. 8.a, 8.b, and 8.c can all be accomplished only using NAND gates in a similar fashion.

9.a

		AB			
		00	01	11	10
C	0	0 1	2 1	6	4 1
	1	1	3 0	7 1	5

$$\text{SOP} = B'C' + A'C' + ABC$$

		AB			
		00	01	11	10
C	0	0	2	6 0	4
	1	1 0	3 0	7	5 0

$$\text{POS} = (B+C')(A+C')(A'+B'+C)$$

9.b

		WX			
		YZ		00	01
00	0	1	4	12	8
01	1	5	13	9	1
11	3	7	15	11	
10	2	6	14	10	1

$$\text{SOP} = \bar{Y}Z + \bar{W}\bar{Y} + \bar{W}X\bar{Z} + W\bar{X}Y\bar{Z}$$

		WX			
		YZ		00	01
00	0	4	12	8	0
01	1	5	13	9	
11	3	7	15	11	0
10	2	6	14	10	

$$\text{POS} = (\bar{Y} + \bar{Z})(W + X + \bar{Y})(\bar{W} + \bar{X} + \bar{Y})(\bar{W} + Y + Z)$$

9.c

		wx			
		00	01	11	10
yz	00	0 1	4 1	12 1	8 1
	01	1	5	13	9 1
	11	3 1	7	15 1	11
	10	2 1	6 1	14 1	10 1

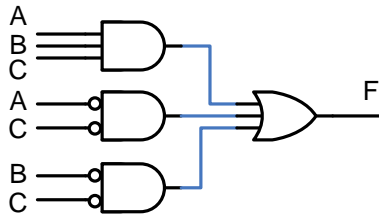
$$\text{SOP} = \bar{z} + \bar{w}\bar{x}y + wxy + w\bar{x}\bar{y}$$

		wx			
		00	01	11	10
yz	00	0	4	12	8
	01	1 0	5 0	13 0	9
	11	3	7 0	15	11 0
	10	2	6	14	10

$$\text{POS} = (W + Y + \bar{Z})(\bar{X} + Y + \bar{Z})(W + \bar{X} + \bar{Z})(\bar{W} + X + \bar{Y} + \bar{Z})$$

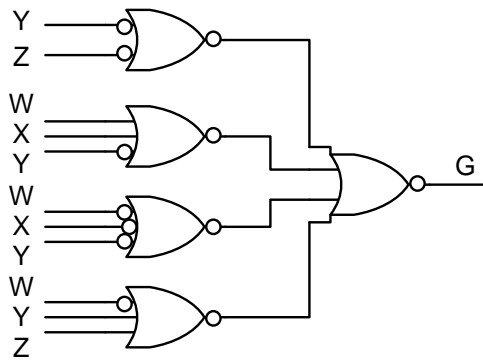
10.a (a) SOP =  $\overline{A}B + \overline{B}C + BC$  shown, but alternately:  $B'C' + A'C' + BC$

- AND-OR implementation



10b.

- NOR-NOR implementation



10c.

- NAND-NAND implementation

