

# CSCI 104

# Linked Lists

CSCI 104 Teaching Team  
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# Lists

- Ordered collection of items, which may contain duplicate values, usually accessed based on their position (index)
  - Ordered = Each item has an index and there is a front and back (start and end)
  - Duplicates allowed (i.e. in a list of integers, the value 0 could appear multiple times)
  - Accessed based on their position ( list[0], list[1], etc. )
- What are the operations you perform on a list?

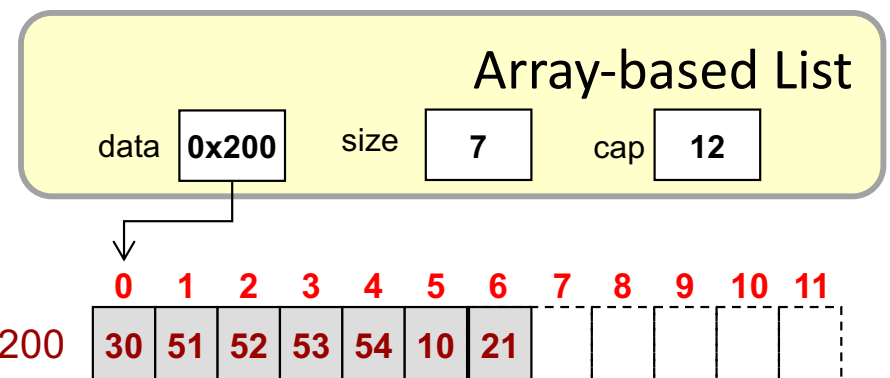
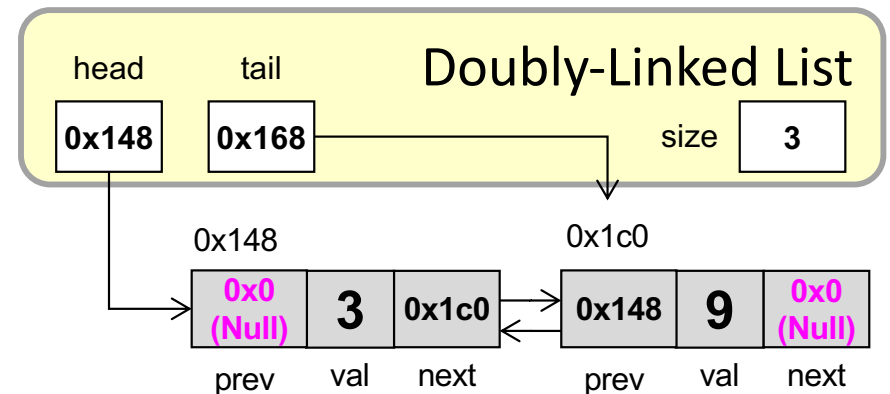
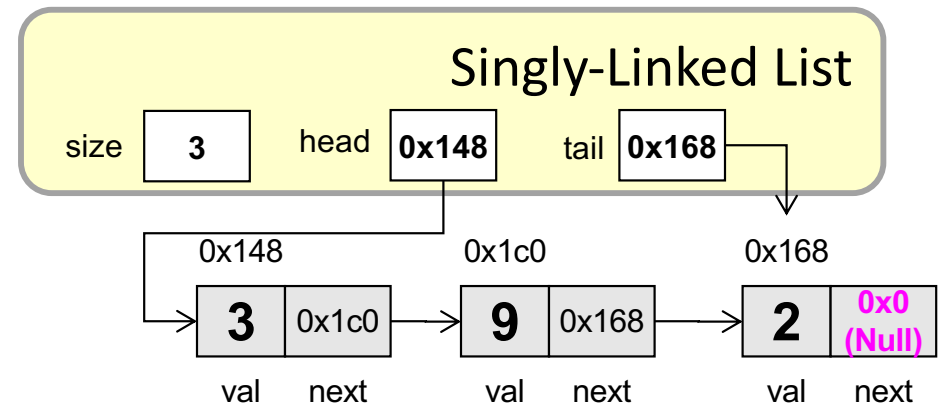


# List Operations

Operation	Description	Input(s)	Output(s)
<b>insert</b>	Add a new value at a particular location shifting others back	Index : int Value	
<b>remove</b>	Remove value at the given location	Index : int	Value at location
<b>get / at</b>	Get value at given location	Index : int	Value at location
<b>set</b>	Changes the value at a given location	Index : int Value	
<b>empty</b>	Returns true if there are no values in the list		bool
<b>size</b>	Returns the number of values in the list		int
<b>push_back / append</b>	Add a new value to the end of the list	Value	
<b>find</b>	Return the location of a given value	Value	Int : Index

# List Implementation Options

- Singly-Linked List
  - With or without tail pointer
- Doubly-Linked List
  - With or without tail pointer
- Array-based List



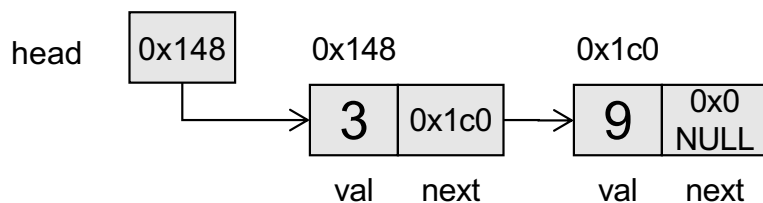
# Implementation Options

## Linked Implementations

- Allocate each item separately
- Random access (get the i-th element) is  $O(\_\_\_)$
- Adding new items never requires others to move
- Memory overhead due to pointers

## Array-based Implementations

- Allocate a block of memory to hold many items
- Random access (get the i-th element) is  $O(\_\_\_)$
- Adding new items may require others to shift positions
- Memory overhead due to potentially larger block of memory with unused locations



# LINKED IMPLEMENTATIONS

# Note

- The basics of linked list implementations was taught in CS 103
  - We assume that you already have basic exposure and practice using a class to implement a linked list
  - We will highlight some of the more important concepts

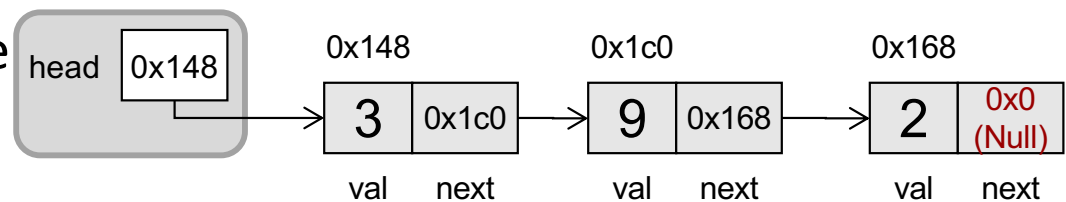
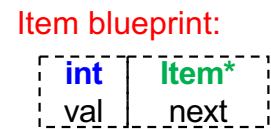
# Linked List

- Use structures/classes and pointers to make 'linked' data structures
- A linked list is...
  - Arbitrarily sized collection of values
  - Can add any number of new values via **dynamic memory allocation**
  - Supports typical List ADT operations:
    - Insert
    - Get
    - Remove
    - Size (*Should we keep a size data member?*)
    - Empty
- Can define a List class to encapsulate the head pointer and operations on the list

```
#include<iostream>
using namespace std;

struct Item {
    int val;
    Item* next;
};

class List
{
public:
    List();
    ~List();
    void push_back(int v); ...
private:
    Item* head_;
};
```

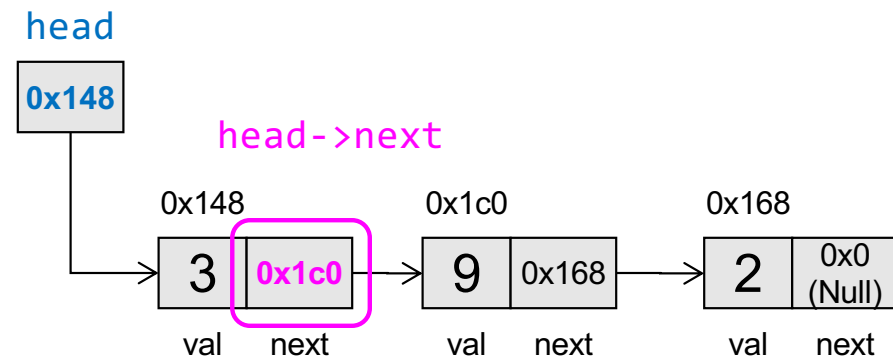


**Rule of thumb:** Still use 'structs' for objects that are purely collections of data and don't really have operations associated with them. Use 'classes' when data does have associated functions/methods.



# A Common Misconception

- Important Note:
  - 'head' is NOT an Item, it is a pointer to the first item
  - Sometimes folks get confused and think head is an item and so to get the location of the first item they write 'head->next'
  - In fact, **head->next** evaluates to the 2<sup>nd</sup> items address



**head->next** yields a pointer to the 2<sup>nd</sup> item!  
**head** yields a pointer to the 1<sup>st</sup> item!

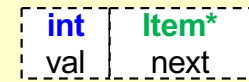
# Don't Need Classes

- Notice the class on the previous slide had *only 1 data member* (the head pointer)
- We don't have to use classes...
  - The class just acts as a wrapper around the head pointer and the operations
  - So while a class is probably the correct way to go in terms of organizing your code, for today we can show you a less modular, procedural approach
- Define functions for each operation and pass it the head pointer as an argument

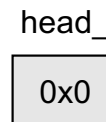
```
#include<iostream>
using namespace std;
struct Item {
    int val;
    Item* next;
};
// Function prototypes
void append(Item*& head, int v);
bool empty(Item* head);
int size(Item* head);

int main()
{
    Item* head1 = NULL;
    Item* head2 = NULL;
    int size1 = size(head1);
    bool empty2 = empty(head2);
    append(head1, 4);
}
```

Item blueprint:



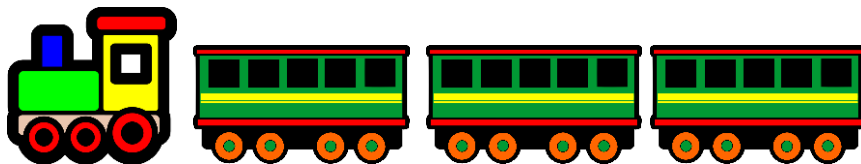
class List:



**Rule of thumb:** Still use 'structs' for objects that are purely collections of data and don't really have operations associated with them. Use 'classes' when data does have associated functions/methods.

# Linked List Implementation

- To maintain a linked list you need only keep one data value: head
  - Like a train engine, we can attach any number of 'cars' to the engine
  - The engine looks different than all the others
    - In our linked list it's just a single pointer to an Item
    - All the cars are Item structs
    - Each car has a hitch for a following car (i.e. next pointer)



Engine =  
"head"

Each car =  
"Item"

```
#include<iostream>

struct Item {
    int val;
    Item* next;
};

void append(Item*& head, int v);

int main()
{
    Item* head1 = NULL;
    Item* head2 = NULL;
}
```

head1

0x0  
NULL

0x0  
NULL

head2

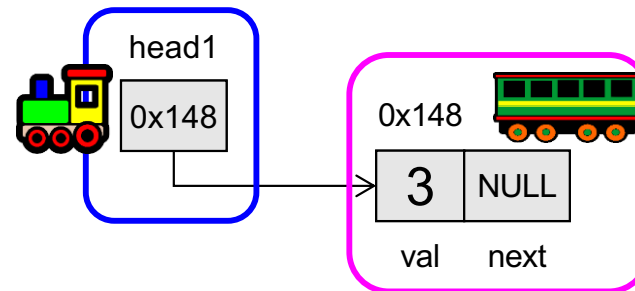
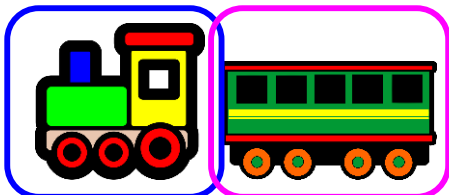
# Append

- Adding an item (train car) to the back can be split into 2 cases:
  - Case 1: Attaching the car to the engine (i.e. the list is empty and we have to change the head pointer)
    - Changing the head pointer is a special case since we must ensure that change propagates to the caller
  - Case 2: Attaching the car to another car (i.e. the list has other Items already) and so we update the next pointer of an Item

```
#include<iostream>
using namespace std;
struct Item {
    int val;
    Item* next;
};

void append(Item*& head, int v)
{
    if(head == NULL){
        head = new Item;
        head->val = v; head->next = NULL;
    }
    else {...}
}

int main()
{
    Item* head1 = NULL;
    Item* head2 = NULL;
    append(head1, 3);
}
```



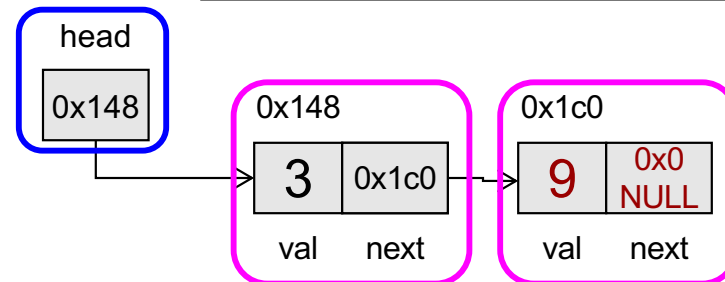
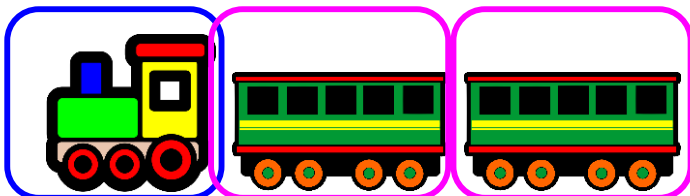
# Linked List

- Adding an item (train car) to the back can be split into 2 cases:
  - Attaching the car to the engine (i.e. the list is empty and we have to change the head pointer)
  - Attaching the car to another car (i.e. the list has other Items already) and so we update the next pointer of an Item

```
#include<iostream>
using namespace std;
struct Item {
    int val;
    Item* next;
};

void append(Item*& head, int v)
{
    if(head == NULL){
        head = new Item;
        head->val = v; head->next = NULL;
    }
    else {...}
}

int main()
{
    Item* head1 = NULL;
    Item* head2 = NULL;
    append(head1,3); append(head1,9);
}
```



# Passing Pointers "by-Value"

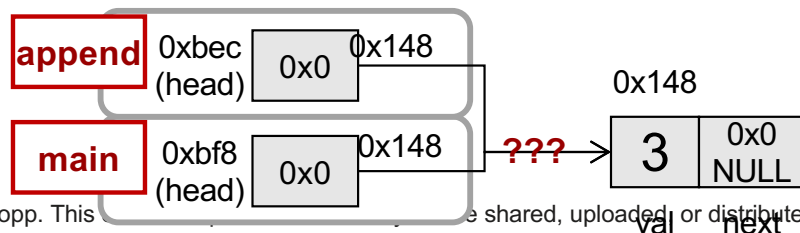
- Look at how the head parameter is passed...Can you explain it?
  - Append() may need to change the value of head and we want that change to be visible back in the caller.
  - Even pointers are passed by value...wait, huh?
  - When one function calls another and passes a pointer, it is the data being pointed to that can be changed by the function and seen by the caller, but the pointer itself is passed by value.
  - You email your friend a URL to a Google doc. The URL is copied when the email is sent but the document being referenced is shared.
  - If we want the pointer to be changed and visible we need to pass the pointer by reference
  - We choose Item\*& but we could also pass an Item\*\*

```
void append(Item*& head, int v)
{
    Item* newptr = new Item;
    newptr->val = v; newptr->next = NULL;

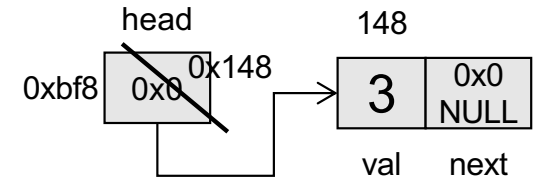
    if(head == NULL){
        head = newptr;
    }
    else {
        Item* temp = head;
        // iterate to the end
        ...
    }
}
```

```
void append(Item** head, int v)
{
    Item* newptr = new Item;
    newptr->val = v; newptr->next = NULL;

    if(*head == NULL){
        *head = newptr;
    }
    else {
        Item* temp = *head;
        // iterate to the end
        ...
    }
}
```



# Passing Pointers by...



```
int main() {
    Item* head1 = 0;
    append(head1, 3);
}
```

**Pointer Passed-by-Value**

```
void append(Item* head, int v)
{
    Item* newptr = new Item;
    newptr->val = v;
    newptr->next = NULL;
    if(head == 0){ head = newptr;}
    else {
        Item* temp = head;
        ...
    }
}
```

```
int main() {
    Item* head1 = 0;
    append(head1, 3);
}
```

**Pointer Passed-by-C++ Reference**

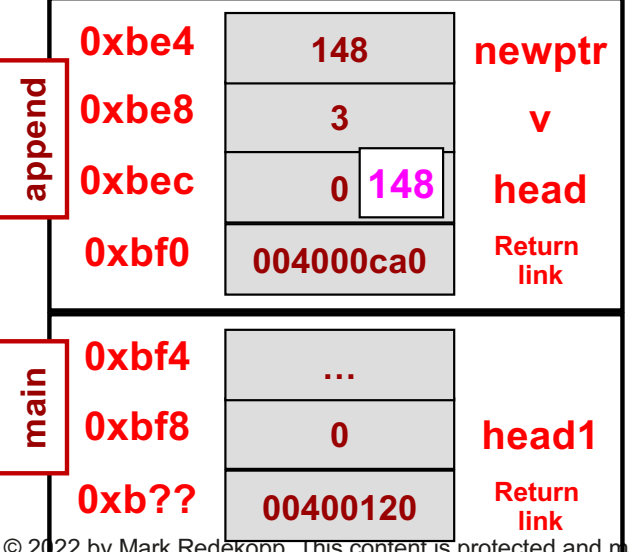
```
void append(Item*& head, int v)
{
    Item* newptr = new Item;
    newptr->val = v;
    newptr->next = NULL;
    if(head == 0){ head = newptr;}
    else {
        Item* temp = head;
        ...
    }
}
```

```
int main() {
    Item* head1 = 0;
    append(&head1, 3);
}
```

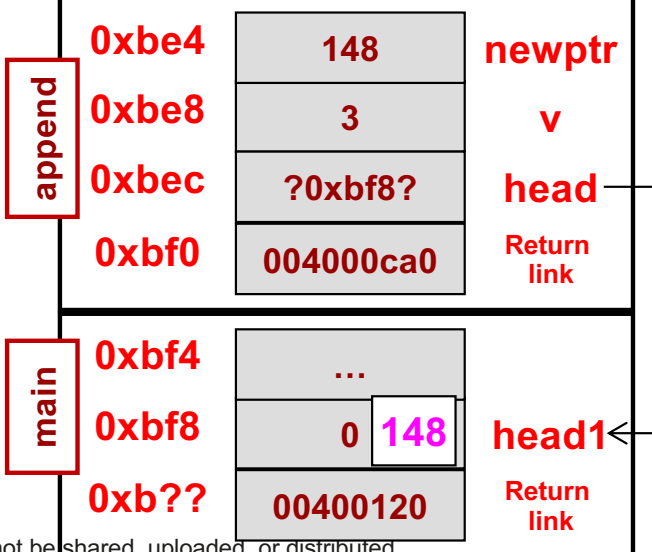
**Pointer Passed-by-Pointer Reference**

```
void append(Item** head, int v)
{
    Item* newptr = new Item;
    newptr->val = v;
    newptr->next = NULL;
    if(*head == 0){ *head = newptr;}
    else {
        Item* temp = *head;
        ...
    }
}
```

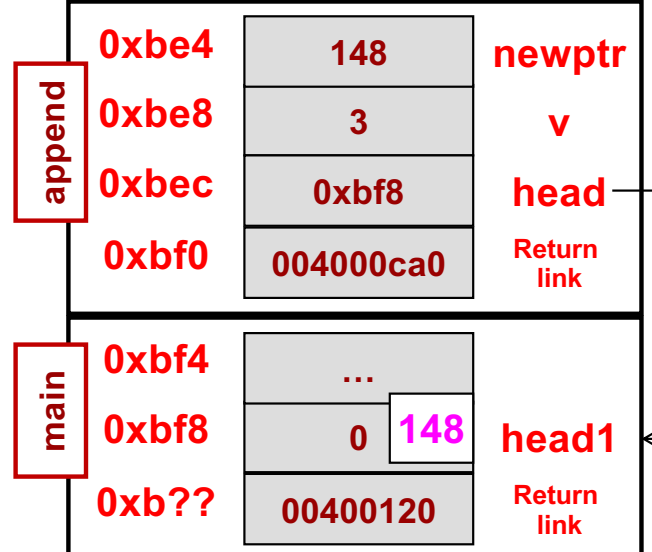
**Stack Area of RAM**



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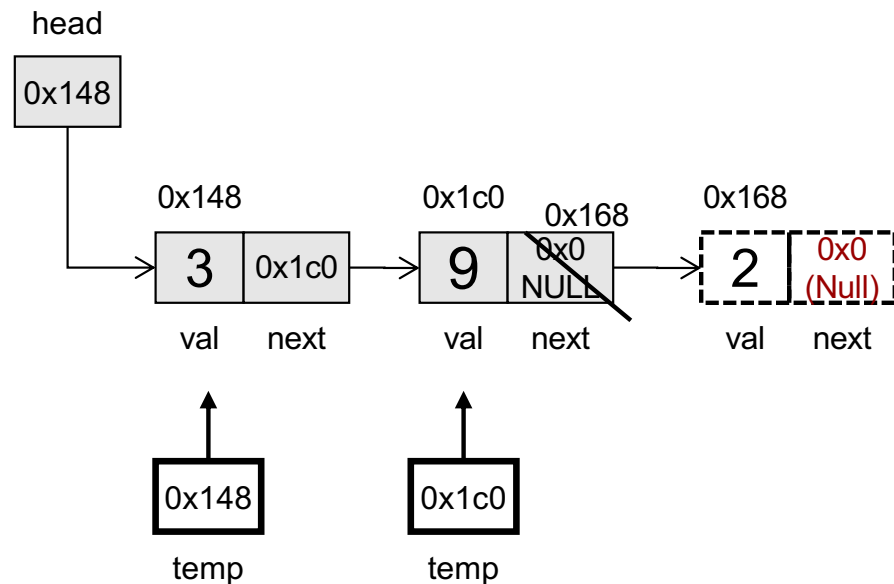


# Iterating Through a Linked List

- Start from head and iterate to end of list
  - Copy head to a temp pointer (because if we modify head we can never recover where the list started)
  - Use temp pointer to iterate through the list until we find the tail (element with next field = NULL)
  - To take a step we use the line: `temp = temp->next;`
  - Optional: Update old tail item to point at new tail item)

```
void append(Item*& head, int v)
{
    Item* newptr = new Item;
    newptr->val = v; newptr->next = NULL;

    if(head == NULL){
        head = newptr;
    }
    else {
        Item* temp = head;
        // iterate to the end
        ...
    }
}
```

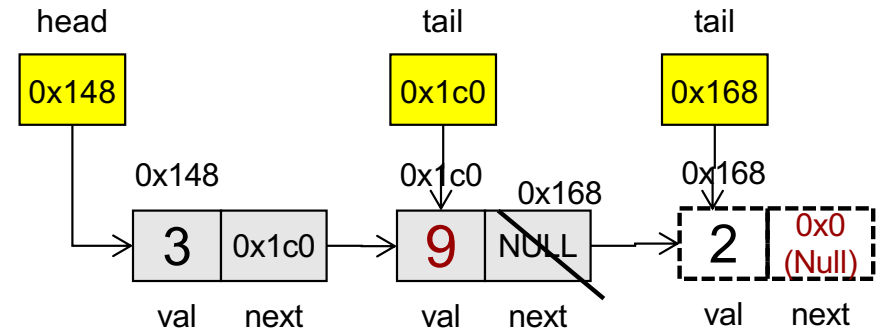


Given only head, we don't know where the list ends so we have to traverse to find it



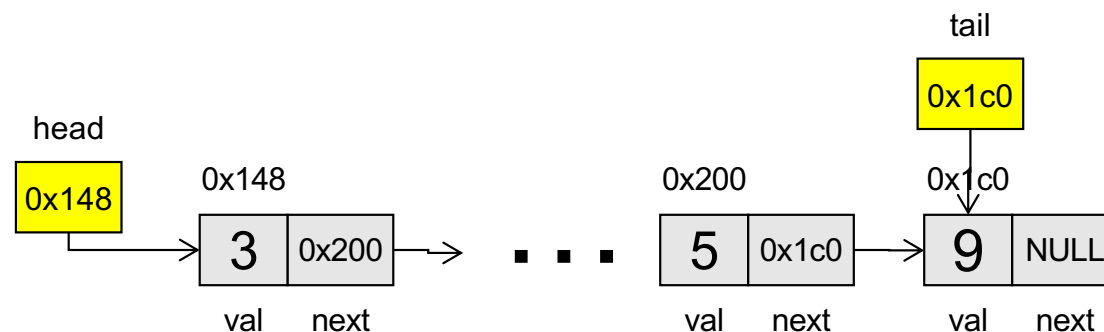
# Adding a Tail Pointer

- If in addition to maintaining a head pointer we can also maintain a tail pointer
- A tail pointer saves us from iterating to the end to add a new item
- Need to update the tail pointer when...
  - We add an item to the end
    - Easy, fast!
  - We remove an item from the end
    - \_\_\_\_\_



# Removal

- To remove the last item, we need to update the 2<sup>nd</sup> to last item (set it's next pointer to NULL)
- We also need to update the tail pointer
- But this would require us to traverse the full list requiring  $O(n)$  time
- **ONE SOLUTION:** doubly-linked list



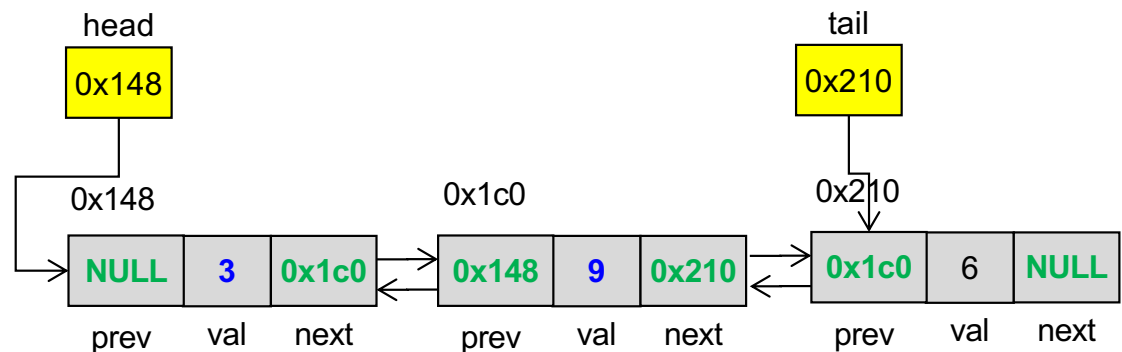
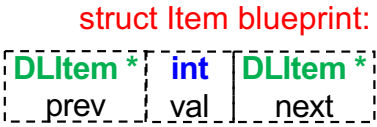
# Doubly-Linked Lists

- Includes a previous pointer in each item so that we can traverse/iterate backwards or forward
- First item's previous field should be NULL
- Last item's next field should be NULL
- The key to performing operations is updating all the appropriate pointers correctly!
  - Let's practice identifying this.
  - We recommend drawing a picture of a sample data structure before coding each operation

```
#include<iostream>

using namespace std;
struct DListItem {
    int val;
    DListItem* prev;
    DListItem* next;
};

int main()
{
    DListItem* head, *tail;
};
```



# Summary of Linked List Implementations

Operation vs Implementation	Push_front	Pop_front	Push_back	Pop_back	Memory Overhead Per Item
Singly linked-list w/ head ptr ONLY					1 pointer (next)
Singly linked-list w/ head and tail ptr					1 pointer (next)
Doubly linked-list w/ head and tail ptr					2 pointers (prev + next)

- What is worst-case runtime of get(i)?
- What is worst-case runtime of insert(i, value)?
- What is worst-case runtime of remove(i)?

# Key Ideas for Linked Lists

- A head pointer is all that is needed to maintain a linked list
- When iterating...
  - Don't lose the head
  - Given a pointer to an item, taking a step to the next node is accomplished with `ptr = ptr->next`
  - Carefully consider when to stop: at the end, one before the end, on the desired item, one before the desired item based on what needs to be updated
- For a singly linked list, use of a tail pointer allows for fast insertion at the end but not removal
- When writing functions that take (head) pointers to linked lists:
  - Always ensure you check and handle if the pointer is NULL
  - If the head/pointer will change, consider how to return that new value (or use pass-by-reference)

# Summary of Linked List Implementations

Operation vs Implementation for Edges	Push_front	Pop_front	Push_back	Pop_back	Memory Overhead Per Item
Singly linked-list w/ head ptr ONLY	$\Theta(1)$	$\Theta(1)$	$\Theta(n)$	$\Theta(n)$	1 pointer (next)
Singly linked-list w/ head and tail ptr	$\Theta(1)$	$\Theta(1)$	$\Theta(1)$	$\Theta(n)$	1 pointer (next)
Doubly linked-list w/ head and tail ptr	$\Theta(1)$	$\Theta(1)$	$\Theta(1)$	$\Theta(1)$	2 pointers (prev + next)

- What is worst-case runtime of `get(i)`?  $\Theta(i)$
- What is worst-case runtime of `insert(i, value)`?  $\Theta(i)$
- What is worst-case runtime of `remove(i)`?  $\Theta(i)$